

# Course 3: Shared Memory

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Course site: *[http://kristaps.bsd.lv/minicourse\\_12\\_2011](http://kristaps.bsd.lv/minicourse_12_2011)*

In prior lectures, we focussed on the environment of high-performance computing:

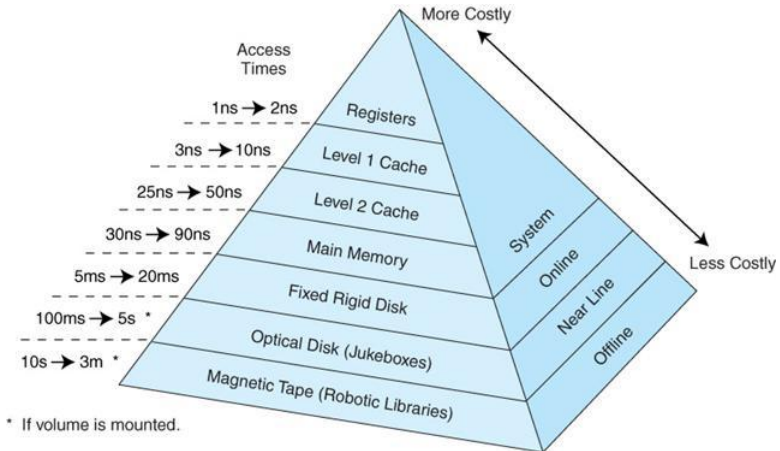
**UNIX** an operating system popular in non-desktop environments

**C** a low-level, minimal programming language

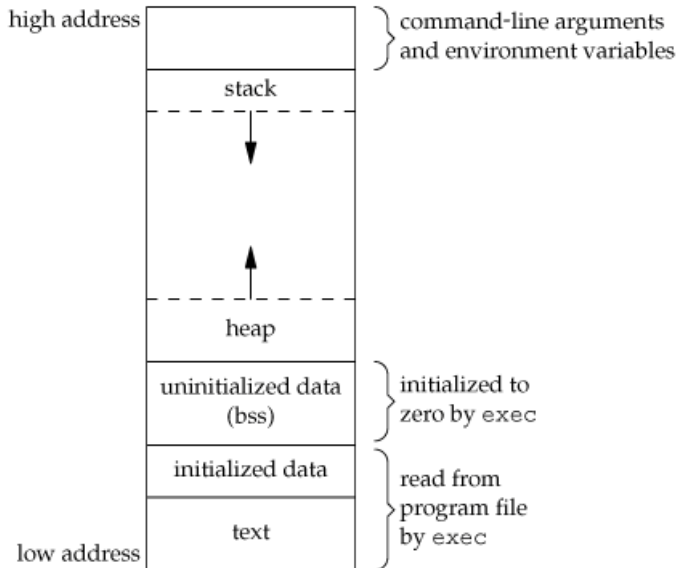
**hardware** purpose-built, high-performance hardware

In this lecture, we consider C programming for high-performance hardware.

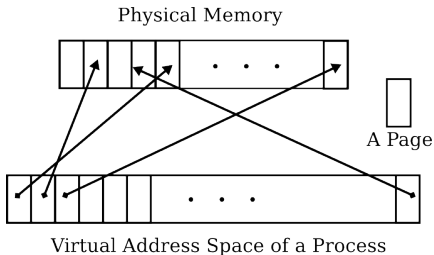
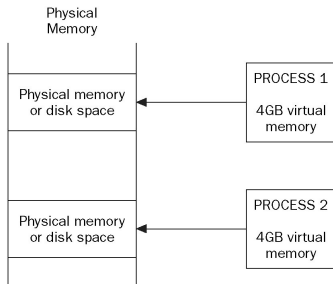
# Review: Hardware



# Review: Process Memory



# Review: Process Memory



# Review: Cache Missing

```
#include <stdio.h>
#include <stdlib.h>

int
main(void)
{
    int                **matrix;
    int                i, j;
    long long int      res;

    matrix = malloc(11000 * sizeof(int *));
    for (i = 0; i < 11000; i++)
        matrix[i] = malloc(11000 * sizeof(int));

    res = 0;
    for (i = 0; i < 11000; i++)
        for (j = 0; j < 11000; j++)
            res += matrix[j][i];

    printf("%lld\n", res);
    return(EXIT.SUCCESS);
}
```

# Review: Cache Missing

D1 misses:	7,584,406	(	7,571,336 rd	+	13,070 wr)
LLd misses:	7,584,016	(	7,570,986 rd	+	13,030 wr)
D1 miss rate:	0.7%	(	0.7%	+	2.9% )
LLd miss rate:	0.7%	(	0.7%	+	2.9% )
LL refs:	7,585,829	(	7,572,759 rd	+	13,070 wr)
LL misses:	7,585,370	(	7,572,340 rd	+	13,030 wr)
LL miss rate:	0.2%	(	0.2%	+	2.9% )

D1 misses:	136,140,031	(	136,126,961 rd	+	13,070 wr)
LLd misses:	121,744,698	(	121,731,668 rd	+	13,030 wr)
D1 miss rate:	14.0%	(	14.0%	+	2.9% )
LLd miss rate:	12.5%	(	12.5%	+	2.9% )
LL refs:	136,141,454	(	136,128,384 rd	+	13,070 wr)
LL misses:	121,746,052	(	121,733,022 rd	+	13,030 wr)
LL miss rate:	4.3%	(	4.3%	+	2.9% )

# Hardware: Stack/Heap Allocation

What are the practical differences between the stack and the heap?

- stack exhaustion
- heap fragmentation
- allocation speed
- overrun protection



# Hardware: Stack/Heap Allocation

```
#include <stdlib.h>

int
main(void)
{
    int b, i, j, *buf, *a = &b;
    for (b = i = 0; i < 1000000; i++) {
        buf = malloc(4096 * sizeof(int));
        for (j = 0; j < 4096; j++)
            buf[j] = (*a)++;
        free(buf);
    }
    return(EXIT_SUCCESS);
}
```

# Hardware: Stack/Heap Allocation

```
#include <stdlib.h>

void
f(int *a)
{
    int i, buf[4096];
    for (i = 0; i < 4096; i++)
        buf[i] = (*a)++;
}

int
main(void)
{
    int a, i;
    for (a = i = 0; i < 1000000; i++)
        f(&a);
    return(EXIT_SUCCESS);
}
```

# Hardware: Stack/Heap Allocation

```
% cc -g -W -Wall -o example6 example6.c
% cc -g -W -Wall -o example5 example5.c
% time ./example5
    0m22.42s real
    0m22.26s user
    0m0.00s system
% time ./example6
    0m27.27s real
    0m27.24s user
    0m0.00s system
```

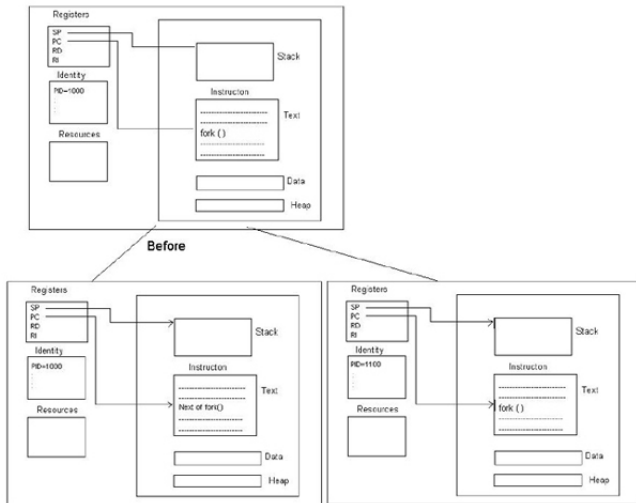
# Shared Memory: Introduction

Consider the dot product:

$$\mathbf{a} \cdot \mathbf{b} = \sum_{i=1}^n a_i b_i$$

In the UNIX and C course, we examined a prime-test using `fork()` (“unshared memory”) to create parallel processes. Can we do the same here?

# Shared Memory: Fork Model



# Shared Memory: Case Study

```
#include <sys/types.h>
#include <sys/wait.h>
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

static float result(void);
static void f(int, float *, float *, int, int);

int
main(void)
{
    float *v1 = NULL, *v2 = NULL, in1, in2;
    size_t sz = 0;

    while (2 == scanf("%f,%f", &in1, &in2)) {
        v1 = realloc(v1, (sz + 1) * sizeof(float));
        v2 = realloc(v2, (sz + 1) * sizeof(float));
        v1[sz] = in1;
        v2[sz++] = in2;
    }

    if (0 == fork())
        f(0, v1, v2, 0, sz / 2);
    if (0 == fork())
        f(1, v1, v2, sz / 2, sz);

    wait(NULL); wait(NULL);
    printf("%g\n", result());
    free(v1); free(v2);
    return(EXIT_SUCCESS);
}
```

# Shared Memory: Case Study

```
static float  
result(void)  
{  
    FILE *f;  
    float in1, in2;  
  
    f = fopen("example3.0.dat", "r");  
    fscanf(f, "%f", &in1);  
    fclose(f);  
  
    f = fopen("example3.1.dat", "r");  
    fscanf(f, "%f", &in2);  
    fclose(f);  
  
    return(in1 + in2);  
}
```

# Shared Memory: Case Study

```
static void
f(int id, float *v1, float *v2, int start, int sz)
{
    FILE *f;
    char buf[64];
    int i;
    float res;

    for (res = 0.0, i = start; i < sz; i++)
        res += v1[i] * v2[i];

    snprintf(buf, sizeof(buf), "example3.%d.dat", id);
    f = fopen(buf, "w");
    fprintf(f, "%g\n", res);
    fclose(f);

    free(v1);
    free(v2);
    exit(EXIT_SUCCESS);
}
```



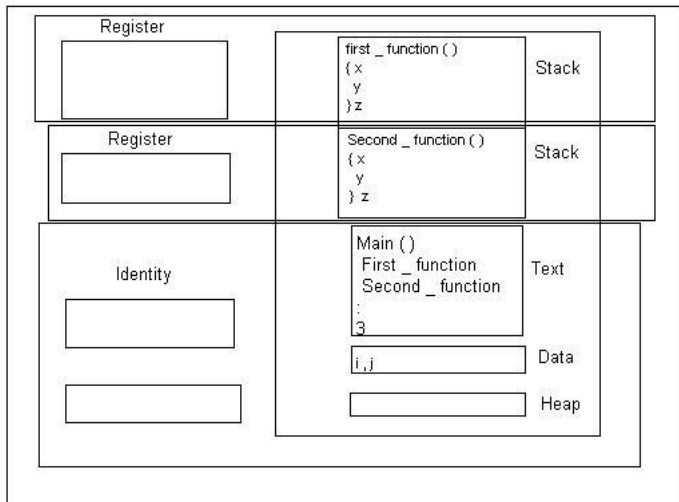
# Shared Memory: Problem

Instead, can we allow multiple processes to share a region of memory? Consider using threads (“light-weight processes”).

- Sharing data section instead of cloning.
- Faster cleanup/teardown.
- Manual synchronisation of memory.

The UNIX standard for threading is *pthread*s, POSIX threads.

# Shared Memory: Thread Model



# Shared Memory: Case Study

```
#include <pthread.h>
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

static float *v1, *v2;
static size_t sz;

static void *
f(void *arg)
{
    int i, start, end;
    float res;

    start = *(float *)arg > 0.0 ? 0 : sz / 2;
    end = *(float *)arg > 0.0 ? sz / 2 : sz;

    for (res = 0.0, i = start; i < end; i++)
        res += v1[i] * v2[i];

    *(float *)arg = res;
    return (NULL);
}
```

# Shared Memory: Case Study

```
int
main(void)
{
    float in1, in2;
    pthread_t p1, p2;

    while (2 == scanf("%f,%f", &in1, &in2)) {
        v1 = realloc(v1, (sz + 1) * sizeof(float));
        v2 = realloc(v2, (sz + 1) * sizeof(float));
        v1[sz] = in1;
        v2[sz++] = in2;
    }

    in1 = -1.0;
    pthread_create(&p1, NULL, f, &in1);
    in2 = 1.0;
    pthread_create(&p2, NULL, f, &in2);

    pthread_join(p1, NULL);
    pthread_join(p2, NULL);

    printf("%g\n", in1 + in2);

    free(v1);
    free(v2);
    return(EXIT_SUCCESS);
}
```

# Shared Memory: Case Study

The code is simpler and scales better. Now consider a simple benchmark:

```
% cc -O2 -g -W -Wall -o example4 example4.c -lpthread
% i=0 ; while [ $i -lt 90000 ] ; do \
> echo $i,$i ; i=$((i+1)) ; done | time ./example4
    0m1.20s real
    0m0.17s user
    0m0.57s system
% cc -O2 -g -W -Wall -o example3 example3.c
% i=0 ; while [ $i -lt 90000 ] ; do \
> echo $i,$i ; i=$((i+1)) ; done | time ./example3
    0m1.52s real
    0m0.15s user
    0m0.56s system
```