

kcgi: sandboxed CGI framework

Or: quantifying the price of web application security

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Part I: i_0 : web application security

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- 1. Client passes HTTP request to web server.
- 2. Web server receives HTTP request from client.
- 3. . . . maps request to CGI.
- 4. ... spawns CGI script.
- 5. ... passes HTTP request to CGI script.
- 6. CGI script processes.
- 7. ... passes HTTP response to server.
- 8. Web server passes HTTP response to client.

6. CGI script processes.

... is really...

6. 6.1 CGI script parses header request information.

6.2 ... reads request into memory/file.

6.3 ... parses key-value pairs from request stream.

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6.4 ... processes key-value pairs.

- 1. Client passes HTTP request to web server.
- 2. Web server receives HTTP request from client.
- 3. ... maps request to CGI.
- 4. ... spawns CGI script.
- 5. ... drop privileges and chroot(2) child.
- 6. ... re-write components of HTTP request into CGI.

- 7. ... passes HTTP request to CGI script.
- 8. CGI script processes.
- 9. ... passes HTTP response to server.
- 10. Web server post-processes HTTP response.
- 11. Web server passes HTTP response to client.



Problem i_0 : adversarial network data directly in contact with system resources.

- Database and application-specific resources.
- System resources (sockets, processes, ...)
- File-system within chroot(2).
- Memory of application (!).
- ▶ ...
- All your base.



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Specifically, problem i_0 is connection of application logic with the code that parses HTTP form data (and HTTP environment) from the CGI request.

Solution $i_0 + 1$?



By splitting apart the parser, we can protect web application logic (the process) from requests exploiting the parser.

We can do better...



1. CGI script spawns child processes before anything else.

- 2. ... hands off standard input to child.
- 3. ... sets up socket pair with child.
- 4. Untrusted child parses header request information.
- 5. ... reads request into memory/file.
- 6. ... parses key-value pairs from request stream.
- 7. ... passes key-value pairs back to CGI script.
- 8. CGI script processes key-value pairs.

The untrusted child can still access system resources.

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We can do better...







Sandboxing is a way of constraining the environment available to a process.

Most well-known sandbox? chroot(2).

More thorough sandboxes: systrace(4) (OpenBSD), Capsicum (FreeBSD), "sandbox" (Darwin), ed(1), ...

By sandboxing the parse sequence, we limit the damage caused by untrusted network data. As for what the application logic does with that data... You're on your own.

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Now on to $i_0 + 2$ and $i_1 \dots$

Part II: $i_0 + 2$: kcgi

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kcgi, kristaps.bsd.lv/kcgi, is a C library that is¹...

- 1. designed to fail
- 2. slow
- 3. resource-intensive

¹It also has manpage documentation, a regression framework, AFL testing framework, automatic HTTP compression, and considerable MIME parsing.

kcgi, kristaps.bsd.lv/kcgi, is a C library that is...

1. designed to fail

1.1 parse and validate *everything* in child process

1.2 sandbox following OpenSSH's example

2. slow

2.1 at least twice as slow as raw parsing

3. resource-intensive

- 3.1 one extra process per CGI client
- 3.2 two extra file descriptors
- 3.3 reads full request into memory twice



Figure : OpenBSD 5.5, nginx, slowcgi(8)



Figure : OpenBSD 5.5, nginx, slowcgi(8)



Figure : Mac OS X "Lion", Apache



Figure : Mac OS X "Lion", Apache

kcgi suffers greatly from the penalty of fork(2), systrace(4), and socketpair(2) I/O. It also suffers from double-allocation of data (original for parse, parsed pairs in parent).

Problem i_1 ? Performance and resource usage.

The allocation problem can be improved by smart programming. But the former?

Solution $i_1 + 1$: FastCGI. This fixes the amount of forking.

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This is a *work in progress*^w^w^w^w.



While this is being finished... *kcgi* is available at kristaps.bsd.lv/kcgi.

Questions?

Thank you!

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